Regular Graphs are Antimagic

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Joint with Yu-Chang Liang and Xuding Zhu Slides available on my webpage

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4	9	2
3	5	7
8	1	6

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62	4	13	51	46	20	29	35
5	59	54	12	21	43	38	28
52	14	3	61	36	30	19	45
11	53	60	6	27	37	44	22
64	2	15	49	48	18	31	33
7	57	56	10	23	41	40	26
50	16	1	63	34	32	17	47
9	55	58	8	25	39	42	24

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62	4	13	51	46	20	29	35		2	59	62	7	18	43	46	23
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52	14	3	61	36	30	19	45	1	58	3	60	17	8	45	22	47
11	53	60	6	27	37	44	22		53	16	5	64	41	20	25	36
64	2	15	49	48	18	31	33		4	57	52	9	32	37	48	21
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16	3	2	13
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9	6	7	12
4	15	14	1

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A Personal Saga: Regular Graphs

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Cor. Just need to check $f(v_1) \neq f(v_2)$ when $v_1, v_2 \in D_i \cup D_{i-1}$.













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Helpful Lemma For each *i*, can partition the edges with distance sum 2i - 1 into paths, so at most one path ends at each vertex. **Pf idea**

Edges induce a bipartite graph; each vertex has degree at most 2.

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