Bootstrap Percolation Thresholds in Plane Tilings using Regular Polygons

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Joint with Neal Bushaw Slides available on my webpage

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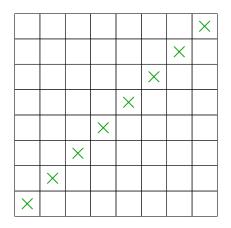
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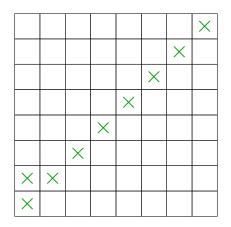
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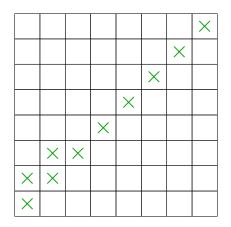
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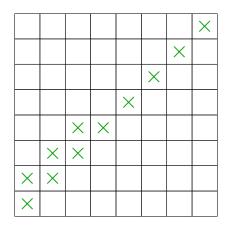
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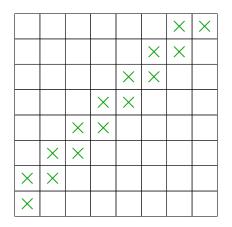
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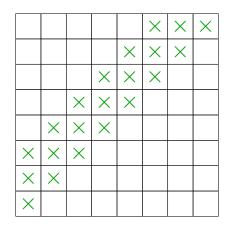
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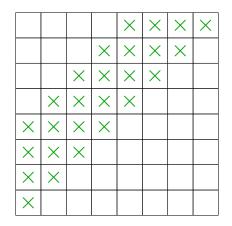
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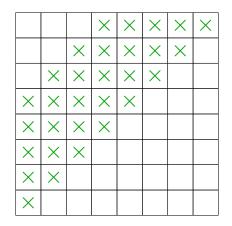
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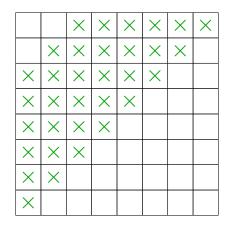
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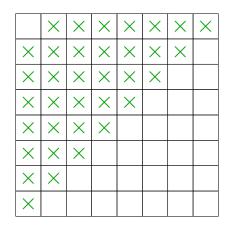
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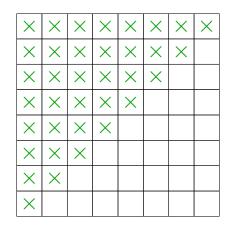
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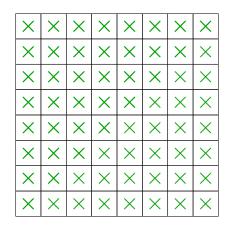
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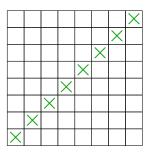
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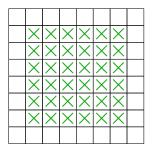


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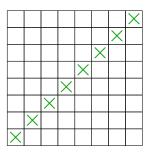
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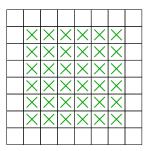




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Ex:





Yes.

No.

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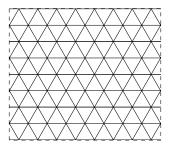
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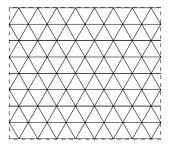
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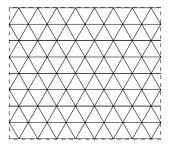
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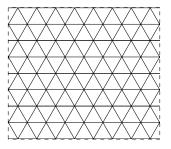




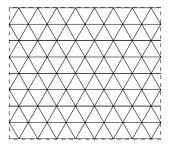
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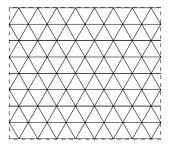
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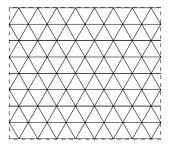
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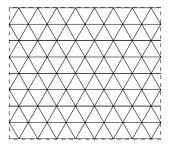
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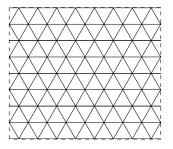
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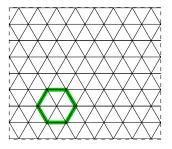
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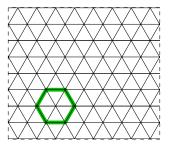
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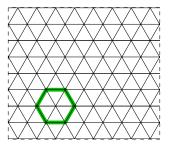
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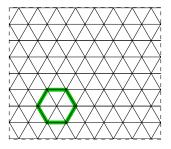
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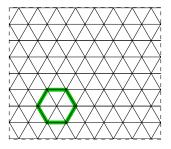
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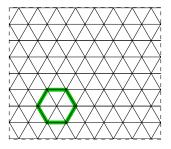
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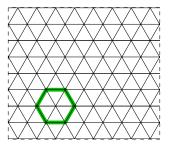
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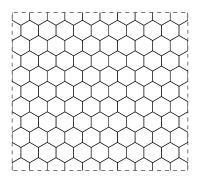


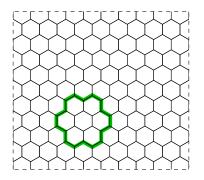
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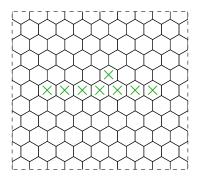
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Cor: The triangular lattice has percolation threshold 1, whenever 0 .

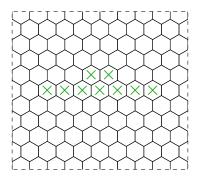




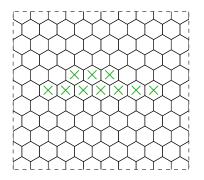
Lem 3: Fix p < 1. For the hex lattice, in the 4-bootstrap model, a *p*-random set \mathcal{I} percolates with prob. 0. **Pf:** Same as Lem 2, but with \bigotimes in place of \bigotimes .



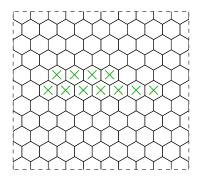
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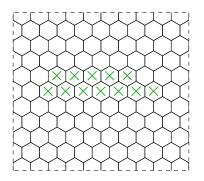
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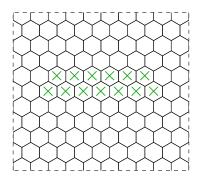
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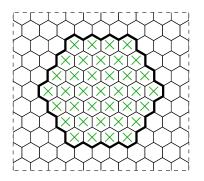
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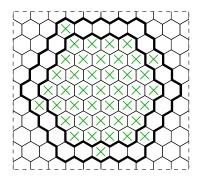
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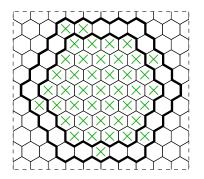
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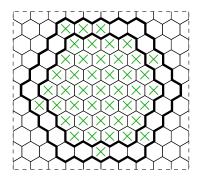
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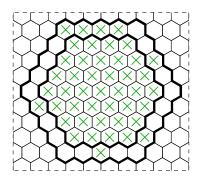
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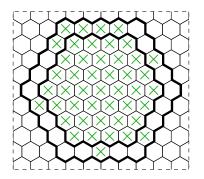
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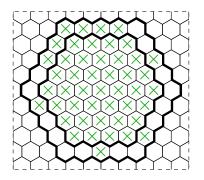
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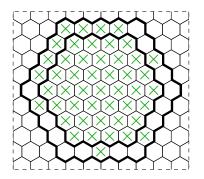
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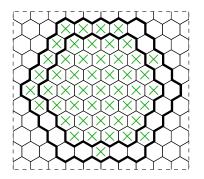
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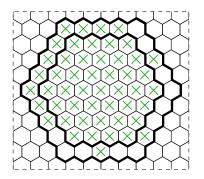
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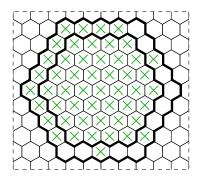
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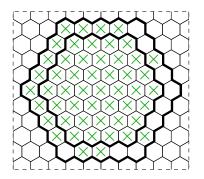
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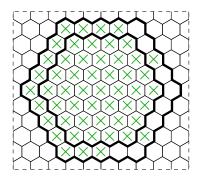
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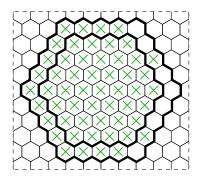
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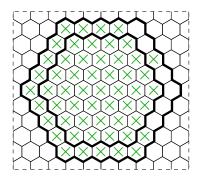
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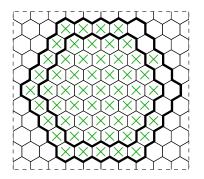
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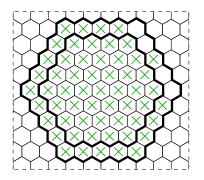
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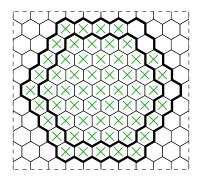
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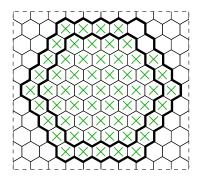
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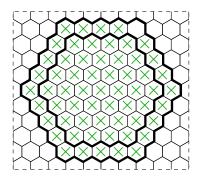
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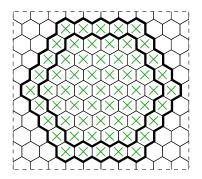
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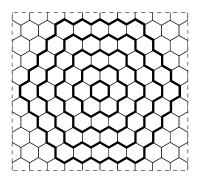


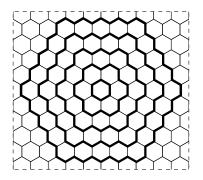
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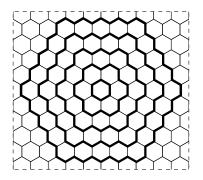
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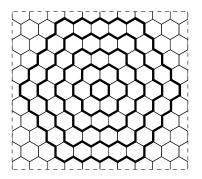




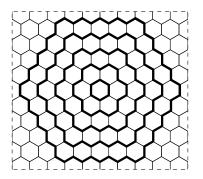
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0-1 Laws

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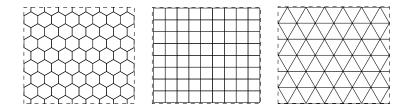
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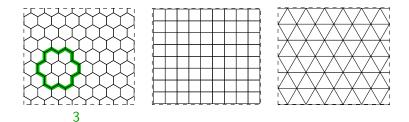
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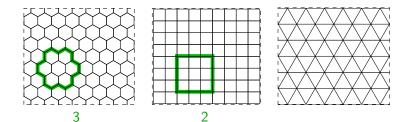
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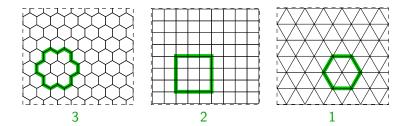
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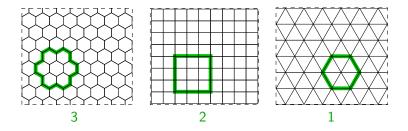
Cor 6: The hex lattice has threshold 3.



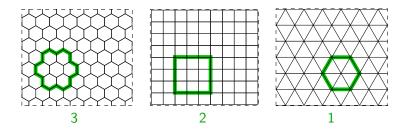






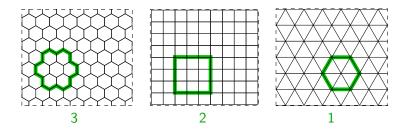


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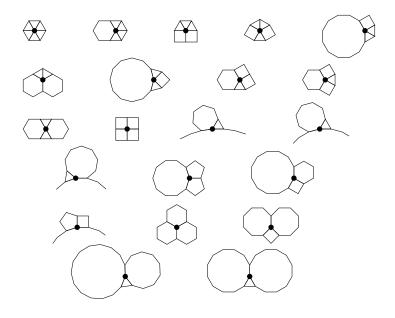


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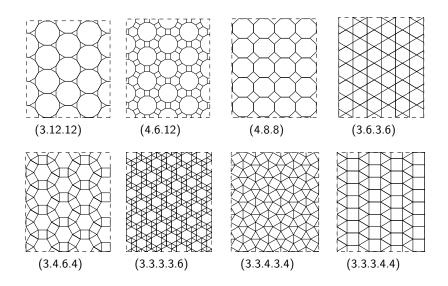
- Allow more face lengths in same graph
- All faces are still regular polygons

How Could Vertices Look?

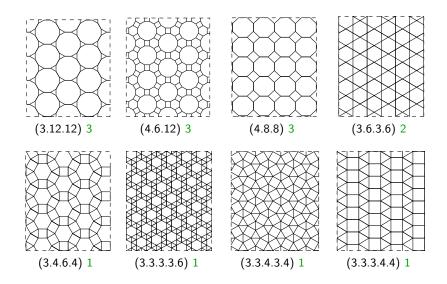
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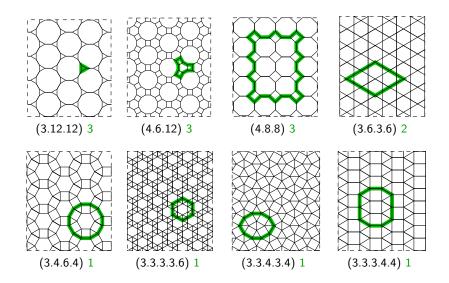
Archimedean Lattices

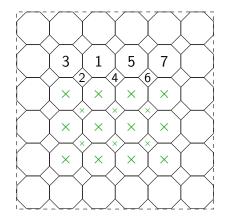


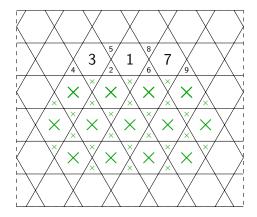
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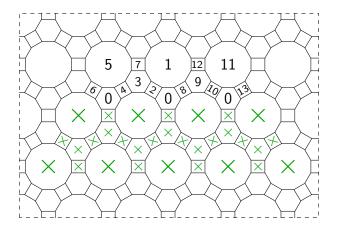


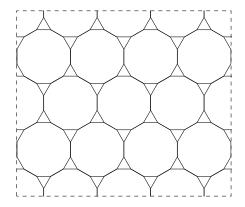
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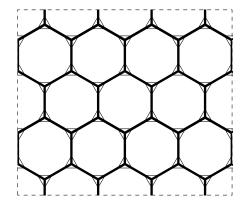












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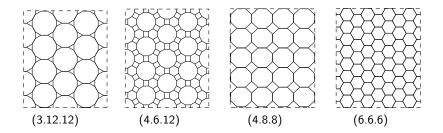
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