An Analogue of Mohar's Conjecture for List-Coloring

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▶ list-assignment L: each vertex v gets allowable colors L(v)

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Main Theorem [C-Mahmoud '23+]: If G is k-regular with $k \ge 3$ and G is connected, then G is k-swappable if $G \ne K_3 \square K_2$.














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- 3. If G is 4-connected and has no 4-wheel, then G is k-swappable (find explicit path between L-colorings).

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- Read more at arXiv:2112.07439